T-110.5150 Applications and Services in Internet

Scaling web applications

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Outline

- What is scalability?
- Scaling strategies
- Scaling network
- Load balancing
- Scaling database
- Monitoring and alerting
- Caching
- Example: eBay architecture (1995-2006)
- CAP (Brewer's) Theorem <- learn more about it!
- Cloud computing

What is scalability?

- A service is called scalable if, adding more resources in a system results in increased performance in a manner proportional to resources added
 - Increased performance? Typically, serving more units of work or handling larger units of work (due to data growth)
- If resources added to improve redundancy of a system: a service is scalable, if adding resources does not result in a loss of performance

What is scalability?

- A scalable system has three characteristics:
 - o It can accommodate an increased usage
 - It can accommodate an increased dataset
 - It is maintainable
- Scalability is not:
 - Performance: high-performing system may not scale

 e.g., it may be fast for 1,000 users and 1 GB data
 but not that fast with 10 times as many users and 10 times data
 - About using any particular technology: Implementation language has little bearing on the scalability of a system
 - Separation of page logic from business logic

What is scalability?

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 - It can accommodate an increased dataset
 - It is maintainable
- Example of a scalable system:
 - <?php sleep(1); echo "Hello world!"; ?>
 - Responds after one second: not fast at all, BUT:
 - Traffic growth is accommodated by adding more web servers: no changes in code
 - Dataset growth is also accommodated: no data stored
 - Code is very maintainable: easy to change
 - Note: coded using not the fastest language

Scaling & hardware

- Hardware is usually expensive at the beginning of any project
 - This is changing now because of cloud computing
- But after project started, cost of software becomes much more expensive
 - Up to a certain point, where, for huge applications, hardware is an issue again
- Meaning that an application is better to be designed and built in such a way to have little or no software work to scale
- Scale by buying more hardware

Strategies for scaling

- Vertical scaling:
 - Get a more powerful version of the same hardware to grow (and throw away the original)
- Horizontal scaling:
 - Get an exact duplicate of the current hardware to grow

Vertical scaling

- Start with a basic setup web server and database server:
 - When each server runs out of capacity, replace it with more powerful
 - When powerful one runs out of capacity, replace it with even more powerful
 - Repeat :)

• Problem:

Reach the limit at some point - price grows exponentially

• But:

- Really easy to design for vertical scaling no changes in software
- Fast alternative If the ceiling for application's usage is known

Horizontal scaling

- Start with a basic setup web server and database server:
 - When each server runs out of capacity, add additional similar
 - When it also runs out of capacity, add additional
 - And so on
- When choosing hardware consider maintenance costs
 - Rack and cable it, install OS, do basic setup, etc.; other issues: space, power, cooling, etc.
- However, at time goes, dealing with hardware additions and failures can become expensive - in other words, increased administration costs:
 - But, systems administration doesn't scale in cost linearly
- Software performance may not scale linearly:
 - Since it needs to aggregate results from all nodes in a cluster, swap message among all of its peers, etc.
 - At some point, it becomes too expensive to add more hardware

Redundancy

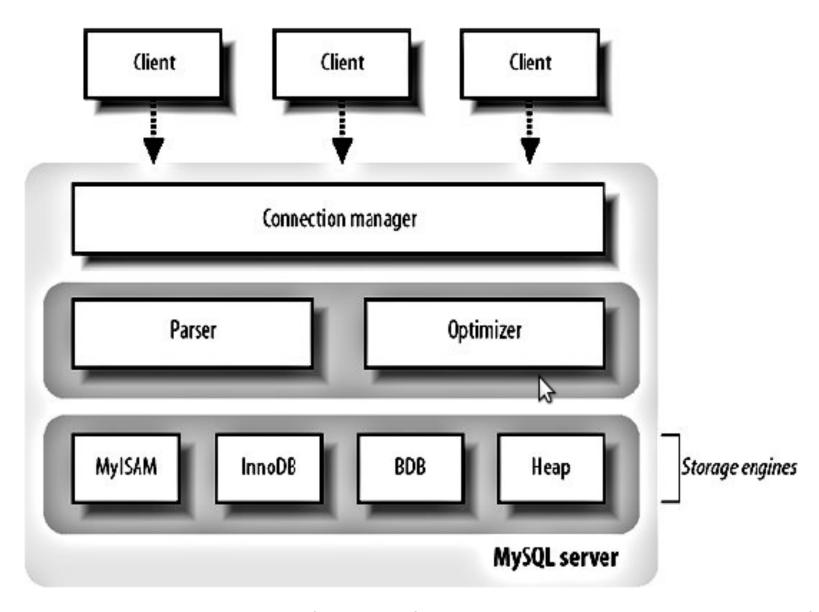
- Machines fail on a regular basis
- One out of every 10,000 machines is expected to die each day
- Have to be prepared to failure of any component
- Spare 'pieces' may be cold, warm and hot:
 - Cold spare: e.g., network switch (physical/software setup and configuration)
 - Warm spare: e.g., queries redirected to a slave database server when a master server dies (configured but needs to be flipped on physically or in software)
 - Hot spare: automatic substitute; e.g., two load balancers are active/passive pair, active takes all traffic and notifies backup balancer via monitoring messages, if active fails, passive stops getting messages and takes over

Scaling network

- As a rule, not a big problem
- Regular networking technologies like gigabit Ethernet provide so much bandwidth that web applications most likely never touch the limits
- Sometimes, an application produces a constant stream of noncritical data and occasional bursts of very important data:
 - May be a problem since Ethernet makes no QoS guarantees
 - Split network into distinct subnets
 - Switches support creation of virtual LANs
- When lots of data has to be transferred between two hosts, high-speed data communication can be used: e. g., InfiniBand (data exchange of up to 100 Gb)

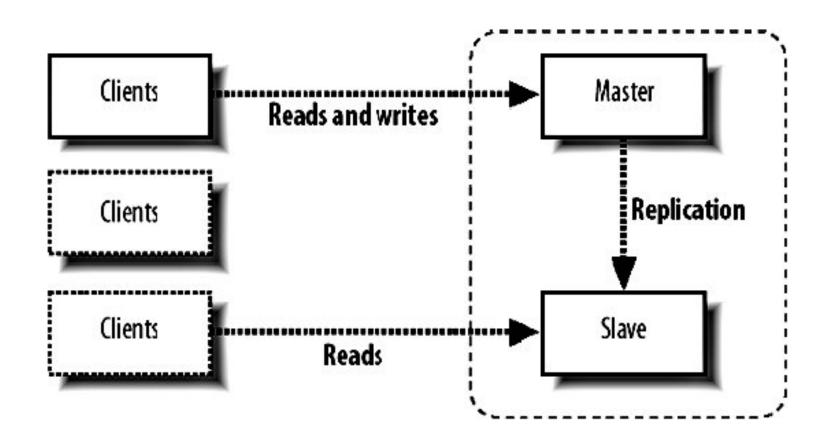
Load balancing

- Vertical scaling: spreading load is the job of operating system scheduler
- Horizontal scaling: there are multiple processors, but no operating system to spread requests between them
 - Several solutions grouped under the term "load balancing"
- DNS-load balancing (dns-issues, hard removal)
- Load balancing with hardware (expensive)
- Load balancing with software

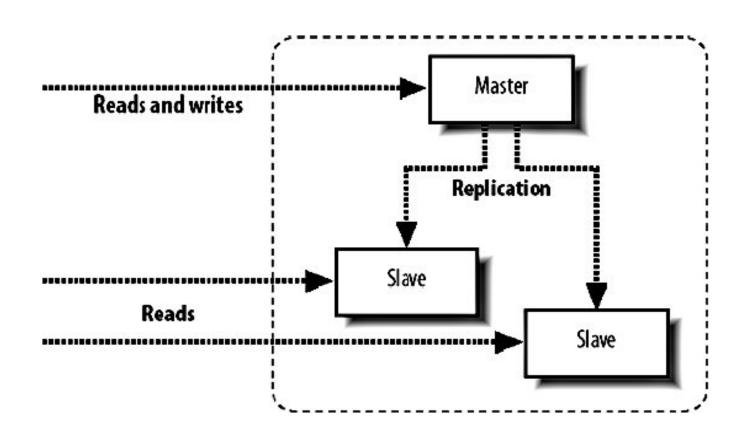


This figure and figures on slides 13-32 have been copied from 'Building scalable web sites', O'Reilly, 978-0596102357, 2006.

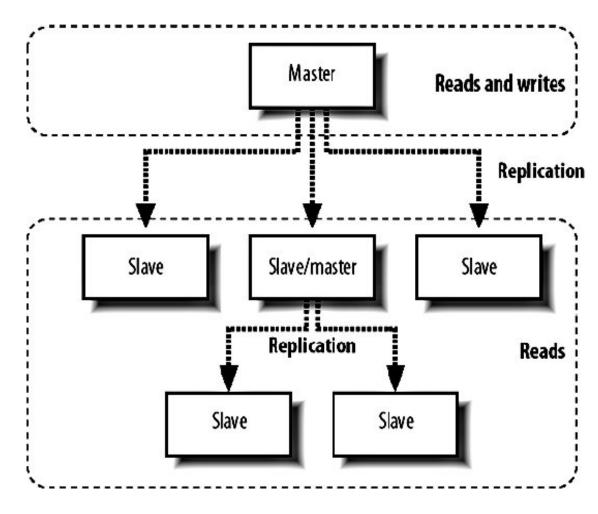
- Web applications typically need a lot more read capacity than write: somewhere between 10 and 100 reads for one write
- Master-Slave replication:



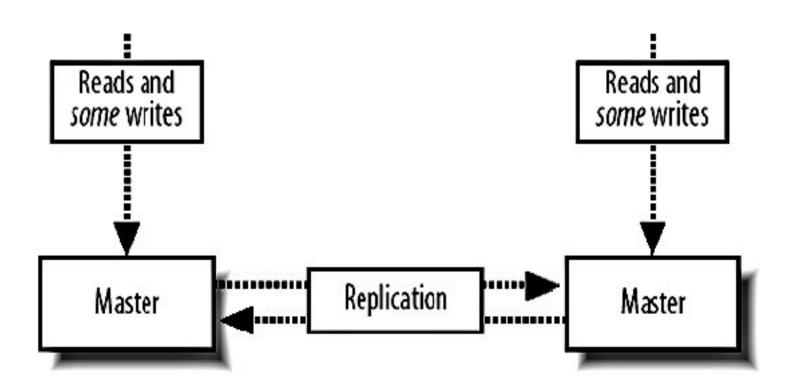
- More read capacity with additional slaves (100 slaves per master in some large applications)
- Slaves aren't guaranteed to be in sync with each other
- Reading scales pretty well, writing is not



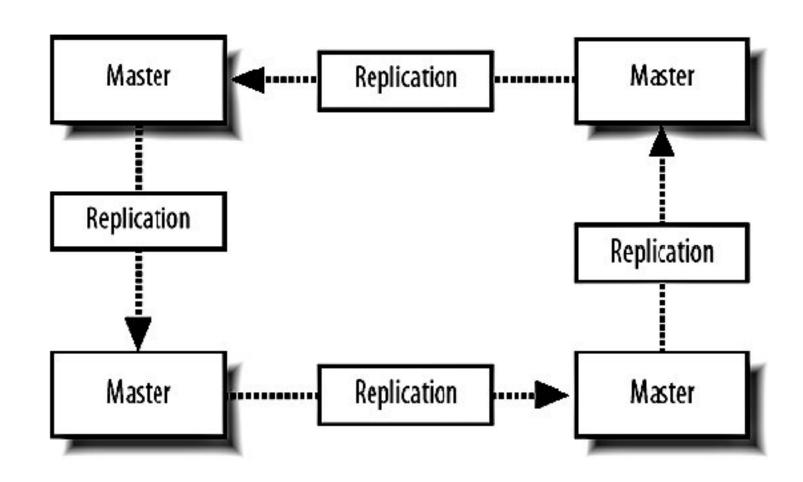
- With lots of slaves, bandwidth required by the master to replicate to all slaves is substantial
- Tree replication



- Master-master replication (each is slave of other)
- Tables with autoincrement primary IDs are problematic

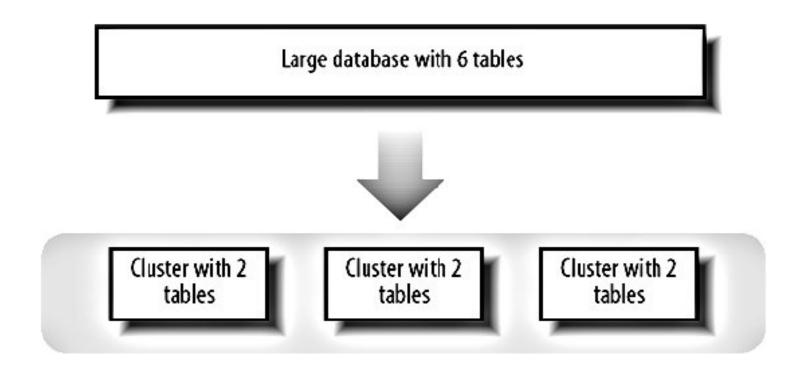


- Masters ring
- At any time, no "true" copy of data
- Failure in a single machine will cause all machines to become stale



- To allow database to scale writes as well as reads, use database partitioning (chop it up into chunks)
- Two ways:
 - Clustering (vertical)
 - Federation (horizontal)

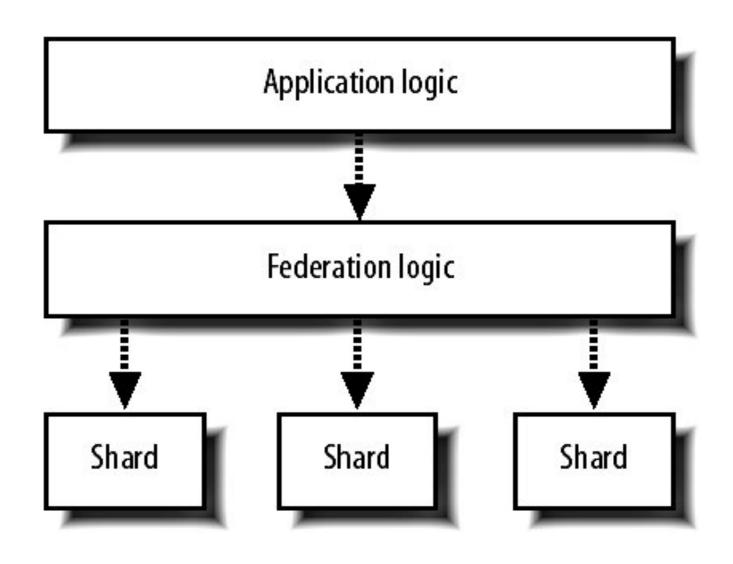
- Clustering:
 - Each cluster contains a subset of tables (joined tables go together)
- Can split only until a single table or a set of joined tables
- Management of cluster is more difficult (different machines have different data)



• Federation:

- Slice the data in the table up into arbitrarily sized chunks
- Chunks of data and the machines that power them are usually referred to as shards or cells
- Avoid cross-shard selects and joins
- The key to avoiding cross-shard queries federate in such as way that all the records you need to fetch together reside on the same shard
- Note: horizontal partitioning in MySQL (5.1 and higher):
 see http://dev.mysql.com/doc/refman/5.1/en/partitioning.html

Federation logic as a separate layer:



Scaling in brief

- When designing a system to be scalable:
 - Design components that can scale linearly by adding more hardware
 - If you can't scale linearly, figure out the return for each piece of hardware added
 - Load balance requests between clusters of components
 - Take into account redundancy
 - Design your components to be fault-tolerant and easy to recover
 - Federate large datasets into fixed-size chunks
- Your application can only scale as well as the worst component in it:
 - Identifying bottlenecks is a must
 - Monitoring infrastructure required

Monitoring

- In a system consisting of many components something always happens
- Do long-term and extensive monitoring:
 - To understand trends and plan for capacity
- Collect and analyze web logs
- Use beacons (a tiny, (usually) invisible image added to the pages of your application for statistical tracking)
- Stream each logfile to a central location

Application monitoring

- Bandwidth monitoring
- Database statistics
- Server runtime statistics (mod_status for Apache)
- Cache statistics

Alerting

- If something crashes or near to crash, you need to know immediately
- That is, there should be a tool that monitors key statistics and alerts whenever some parameter goes below or above a certain value
- Uptime checks (to check that a service or component is up)
- Threshold checks
- Low-watermark checks

Identifying bottlenecks

- As Donald Knuth said: premature optimization is the root of all evil (or at least most of it) in programming
- Optimizing any part of your application before finding out whether that component requires optimization is a waste of time

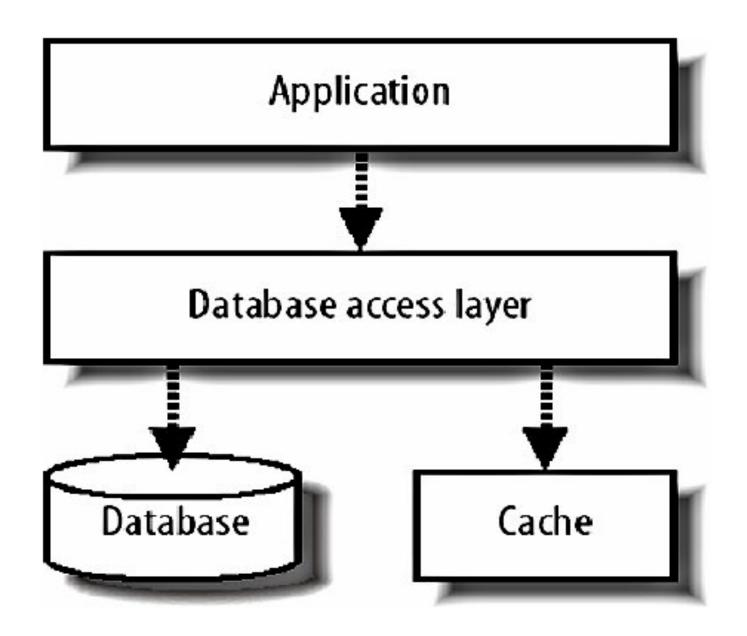
Database throughput

- In many web applications, the biggest bottleneck is database throughput, usually caused by disk I/O
- When a database being a bottleneck, it is generally about the time between a query reaching the database server and its response being sent out
- Certain kinds of data are good for caching:
 - Get set very infrequently but fetched very often
 - E.g., account data: you might want to load account data for every page that you display, but you might only edit account data during 1 in every 100 requests

Adding caching

- When fetching an object, we first check in the cache
 - If object exists, read it from the cache and return it to the requesting process
 - If it doesn't exist in cache, we read the object from the database, store it in cache, and return it to the requester
- When changing an object in the database, either for an update or a delete, it needs to be invalidated in cache

Adding caching



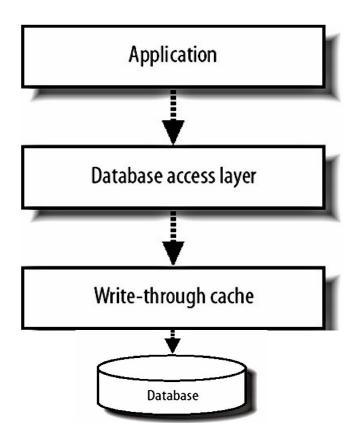
Caching: memcached

- memcached: memory cache demon
- Open source memory-caching system designed to be used in web applications to remove load from database servers
- Supports up to 2 GB of cache per instance
 - But allows you to run multiple instances per machine and spread instances across multiple machines
- Native APIs available for PHP, Perl and other common languages

Caching as entire level

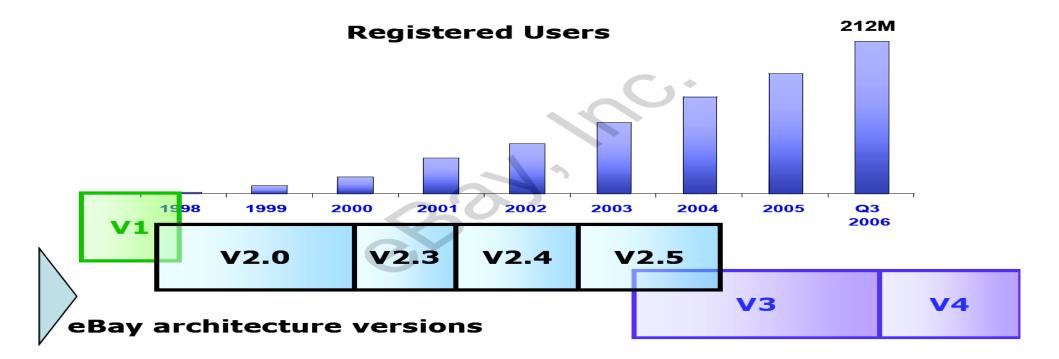
• Problems:

- Reads are caching, while all writes need to be written to database synchronously and cache either purged or updated
- Cache-purging and updating logic gets tied into our application and adds complexity
- Cache as layer of its own:



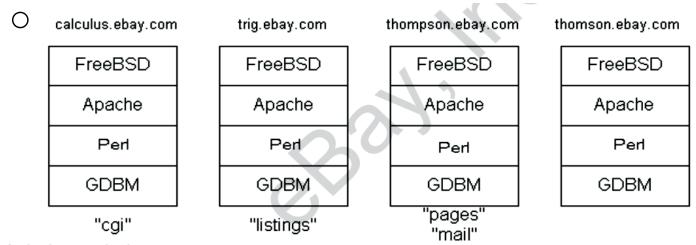
Example: eBay architecture (1995-2006)

- See http://highscalability.com/blog/2008/5/27/ebayarchitecture.html (figures taken from the eBay architecture talk, http://www.addsimplicity.com/downloads/eBaySDForum2006-11-29.pdf)
- 2006 eBay figures:
 - Over 200mln registered users
 - Over one billion photos
 - Over two petabytes of data
 - Over 26 billions executions per day

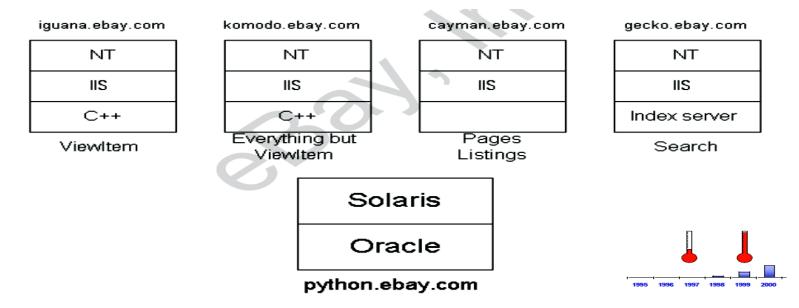


Example: eBay architecture

 1995-97 (v1; every item was a separate file generated by perl script):

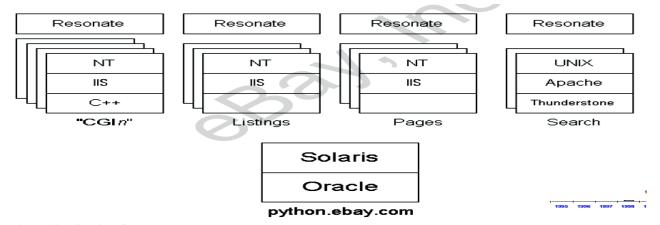


• 1997-99 (v2; C++, MS, Oracle):

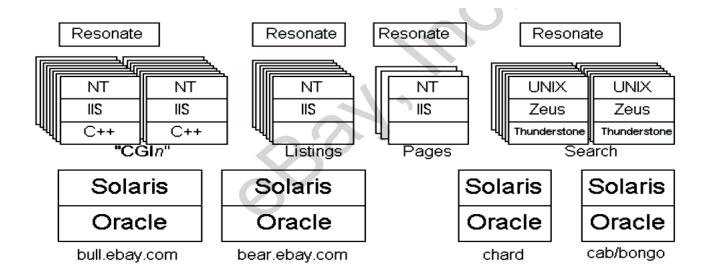


Example: eBay architecture

• 1999 (v2.1; server grouped into pools, front-end balancing and failover, db scaled vertically to a larger machine):



• 1999-2001 (v2.3-2.4; horizontal scaling of servers continued, second db for failover, db 'splitting' started):



Example: eBay architecture

- Observation:
 - Change of technology stack all the time
 - 2002: moved from C++ to Java
 - Later: Oracle->MySQL (sharding)
- Scaling data tier
 - Segmentation by function
 - User hosts, item hosts, account hosts, feedback hosts, ...
 - More that 70 other
 - Horizontal splits within function
 - Introducing logical database hosts (additional layer)

CAP Theorem

- At PODC'00 (in an invited talk "Towards Robust Distributed Systems"), Brewer made the conjecture:
 - It is impossible for a web service to provide the following three guarantees:
 - **■** Consistency
 - **A**vailability
 - **■** Partition-tolerance
 - All three are expected and very desirable from real-life web services
- In 2002, Gilbert and Lynch of MIT formally proved the Brewer's conjecture
 - Prove is rather simple

CAP Theorem

- Consistency (all nodes see the same data)
 - Formally, atomic consistency
 - Requests to a distributed shared memory are executed like they are executing on a single node (i.e., one by one)
 - Property: read operation that begins after write operation completes must return the value/result of this write operation
- Availability (node failures do not prevent others to operate)
 - Every request received by non-failing node must result in a response (or every request must terminate)
- Partition-tolerance (loss of some messages doesn't disrupt the system operation)
 - A network is partitioned if all messages sent from nodes in one component of the partition to nodes of other partition are lost

CAP Theorem

- No Consistency
 - web caching, DNS, NoSQL dbs
- No Availability
 - distributed databases, majority protocols
- No partition-tolerance
 - single-site databases, LDAP

Cloud computing

- Servers are not used all the time:
 - Need to have them for redundancy
 - Idle time is cost-inefficient
- Computational power of 1000 servers for one hour costs the same as using power of one server for 1000 hours:
 - Cloud computing services such as Amazon Elastic Compute Cloud made the first option (1000 servers for 1 hour) available
 - One can start with one virtual instance, add more almost instantly when needed (if, of course, application was designed correspondingly)

Literature

- Books&articles:
 - o Building scalable web sites, O'Reilly, 2006
 - Developing large web applications, O'Reilly, 2010
 - o The art of capacity planning, O'Reilly, 2008
 - Performance by design: Computer capacity planning by example, Prentice Hall, 2004
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